Codeforces Review

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.1 2013D - 1900

Problem 1. ¹ Given an array A and an operator p that for each a_i and a_{i+1} where $i \ge 1$, p may decrease $|a_i - a_{i+1}|$. p can be performed by infinite times and find the minimum $|max_{i=1}^n(a_i) - min_{i=1}^n(a_i)|$

Solution 1. Suppose b_i is the result of performing p on a_i given the prefix of array $A \{a_1, a_2, \ldots, a_i\}$. Create a stack S to load these b_i and the count of it c_i . Then $\{a_1, a_2, \ldots, a_i\}$ is stored in S as $(b_1, c_1), (b_2, c_2), \ldots, (b_m, c_m)$. And we keep the pairs in ascending order of b_i . For each new a_{i+1} , merge it to the top if $a_{i+1} < b_m$ and then merge the top downwards until $b_k > b_{k-1}$. This way, each a_i is loaded in the array for once, the time of merging until a_i is at most i, so the time complexity if O(n).

.2 2014H - 1900

Problem 2. ² Given an array A, check for each a_i , if there exists $A_m = \{a_{m_1}, a_{m_2}, \ldots, a_{m_j}\}$ s.t. $a_i \in A_m$ and j is even.

Solution 2. If the size of A is small, for instance, 166, then first hashing the A in a much bigger set, for instance, $\{1, 2, ..., 2^{64}\}$, and check if the xor sum of the array is 0. The hash is very important, without which, several different numbers can also get xor sum 0 (e.g., $\{1, 2, 3\}$). The possibility of reaching such bad situation after hashing is $\frac{1}{2^{64}}$.

.3 2035D - 1800

Problem 3. ³ Given an array A and an operator p that for each a_i and a_j that i < j, p can update a_i to $a_i \gg 1$ and a_j to $a_j \ll 1$. p can

¹https://codeforces.com/problemset/problem/2013/D

²https://codeforces.com/problemset/problem/2014/H

³https://codeforces.com/problemset/problem/2035/D

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be performed by infinite times and find the maximum sum of all the prefixes of A.

Solution 3. A little similar to 2013D. In 2013D, we store b_i and c_i , the information of a_i after p in a stack S. This approach is appliable in this problem too. In this problem $b_i = \min\{b_i^j | b_i^j \ll c = a_i\}$, and c_i is the sum of all c_k that k < i that can reach the maximum prefix sum from 1 to i. This sum is obtained by merging from the top of the stack downwards. If $c_k(k < i)$ is added to c_i , then b_k is popped out from the stack and added to the final sum. After the merging is terminated, we push (b_i, c_i) onto S. Each (b_i, c_i) pair is at most pushed to S by once and popped by once. So the time complexity is O(n)

.4 2009G1-1900

Problem 4. ⁴ Given an array A and an operator op that can change any element a_i to any another value. Find the minimum step required to make A a consecutive array in which $a_{i+1} - a_i = 1$ for any *i*.

Solution 4. Create an array B that $b_i = a_i - i$. Count the frequency f_i of each b_i and the answer is $A.size() - max(f_i)$.

.5 2002D1, 2021C2 - 1900

Problem 5. ⁵⁶ Continuously query about a property *P* with an update before each query.

Solution 5. Find another easy-to-maintain property P' s.t.if P' is satisfied then P is satisfied. The time complexity of checking for P' after each update is O(1) or O(logn).

⁴https://codeforces.com/problemset/problem/2009/GI

⁵https://codeforces.com/problemset/problem/2002/DI

⁶https://codeforces.com/problemset/problem/2021/C2

.6 1991E - 1900

Takeaway 1. ⁷ If I want to use two colors to color an undirected graph G, in which for every v_1 , v_2 connected by edge E, their colors are different, the only thing to judge if I can make the coloring is check if the graph is bipartite.

.7 1991D - 1900

Takeaway 2. ⁸ If a - b = kx, and x is the power of 2, then $a \bigoplus b = x$ where \bigoplus means XOR

.8 2039C2 - 1800

Takeaway 3. ⁹ Some basic math that is useful for this kind of CF problems. Suppose $a \ge b$, then $a - b \le a \bigoplus b \le a + b \le 2a$ since \bigoplus is just addition without carry or subtraction without borrowing. Suppose a > b, then $lcm(a,b) \le 2max(a,b) \le 2a$.

⁷https://codeforces.com/problemset/problem/1991/E

⁸https://codeforces.com/problemset/problem/1991/D

⁹https://codeforces.com/problemset/problem/2039/C2